



(My) Uncovering Exploitable Firmware Internship

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Agenda

- About Me
- UEFI
 - Background
 - A brief history
 - Supply chains
 - Vulnerabilities
- HARDEN
 - Detecting vulnerabilities
 - Scaling up detection

About Me

- Originally from Allentown, PA
- Currently: University of Maryland at College Park
 - Studying computer science
 - Minor in robotics and autonomous systems
 - Planning to graduate in Dec. 2023
- UMDLoop
 - Avionics Systems Lead
 - Telemetry for a tunnel boring machine
 - Software and hardware for a Mars rover



Background on UEFI

A Brief History

- PCs use firmware to initialize hardware/software
 - Historically: [Legacy] BIOS (c. 1981)
 - Limited functionality
 - Address-constrained
 - No support for fancy features
 - Network boot?
 - Development was difficult + highly machine specific
 - Today: UEFI (c. 2005)
 - **U**niversal **E**xtensible **F**irmware **I**nterface



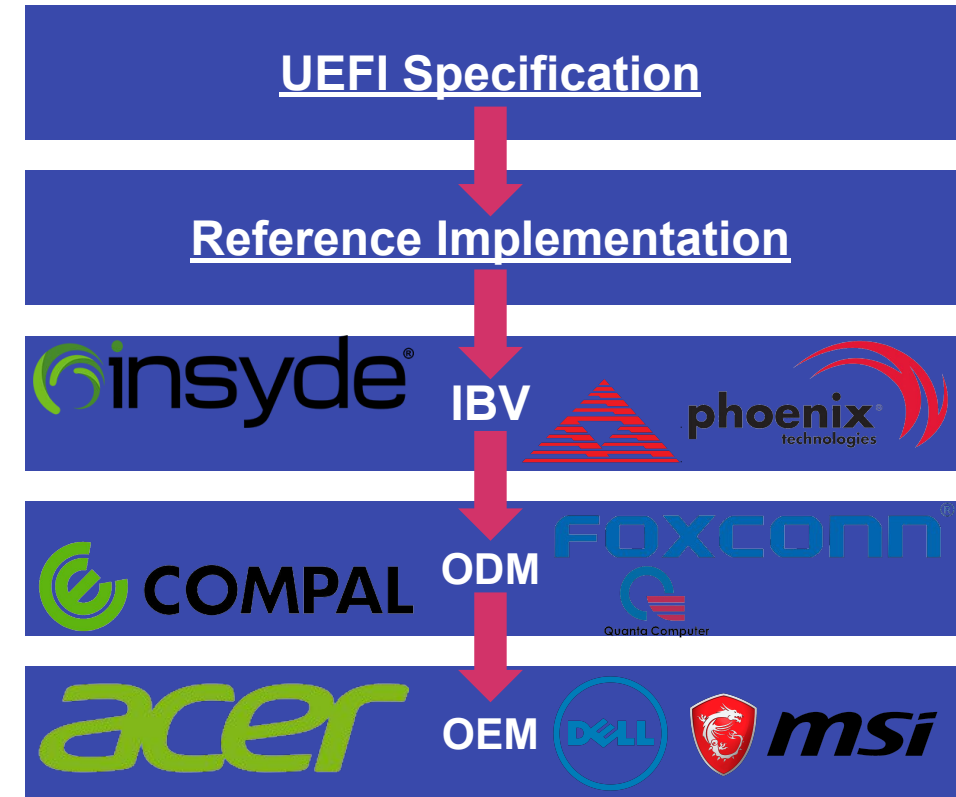
UEFI

- First developed by Intel (2005)
- A common standard
 - Vendors implement independently
 - Reference implementation: Tianocore EDK II
- Not limited to just boot-time services
 - Manages runtime kernel - hardware interaction
 - Most UEFI drivers run in driver execution environment (DXE)
 - Equivalent to ring 0/kernel mode



UEFI Supply Chains

- Lots of different companies/organizations involved!
 - Bug fixes take a *long* time to reach end-users
 - Independent BIOS Vendors
 - Original Design Manufacturers
 - Original Equipment Manufacturers
- How UEFI firmware is developed/distributed is dependent on all of these organizations
 - Annoying packaging practices (Dell 😡)



Taken from Rylan's JTB on UEFI (thanks Rylan)



UEFI Vulnerabilities

- Types
 - Double GetVar
 - GetSet
 - SMM Callouts
 - SMM CommBuffer poisoned pointers
- In general
 - Vulnerabilities are simple, but hard to find because of the supply chain

Double GetVar & GetSet

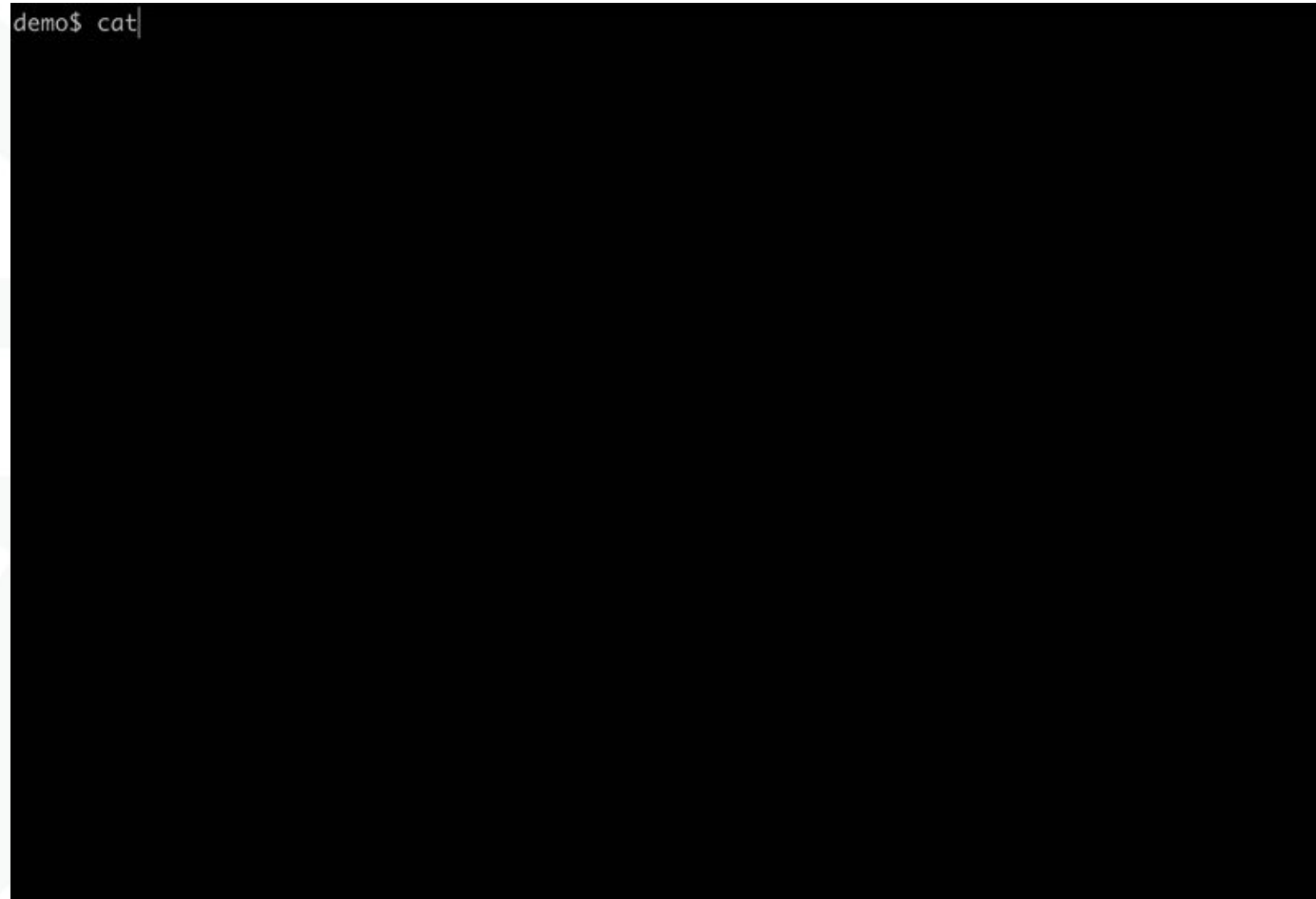
- Runtime service that gets a value from NVRAM
 - NVRAM - non-volatile RAM
 - Persistent key/value storage for UEFI variables
- Takes a *DataSize* pointer
 - Input: Size of the data buffer we are writing the variable value to
 - Output: # of bytes that the variable occupies
 - Can be longer than the buffer
 - Returns an error if it is
- Two GetVariable calls without sanitization? Potential overflow!

GetSet

- Very similar to Double GetVar
 - SetVariable
 - Sets a value in NVRAM
 - Consecutive calls to GetVariable and SetVariable without sanitization
 - Exposes EFI variables

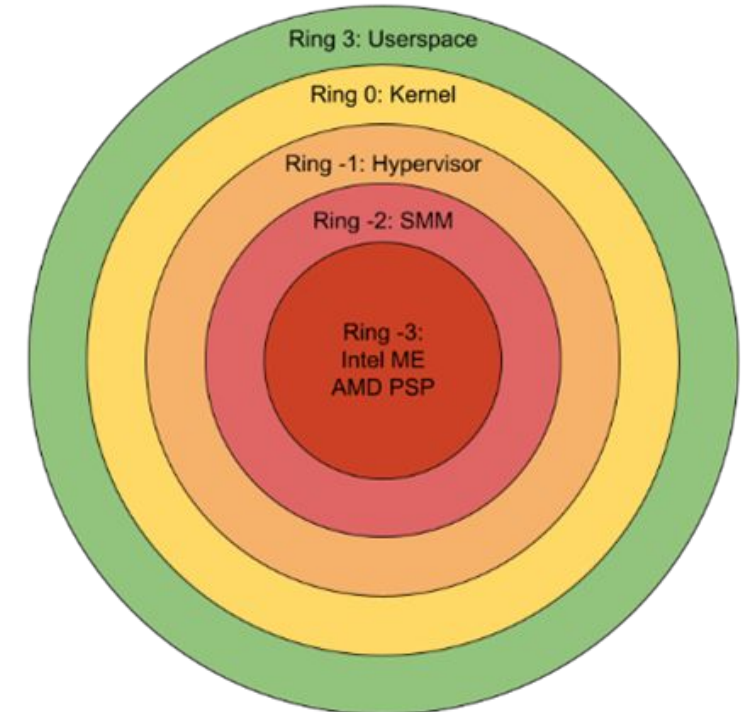
Double GetVar & GetSet (cont.)

```
demo$ cat|
```



System Management Mode

- x86 processors have an execution mode called System Management Mode (SMM)
 - Runs highly privileged code
 - Can interact with all physical memory (even though it shouldn't)
 - Will run above hypervisors as well
 - Equivalent to ring -2
 - ring 0: kernel space
 - Mostly invisible to the OS
 - Invoked from the kernel, or from a hardware interrupt
- Can interact with SPI flash (and all other hardware)
 - Install rootkits that can persist *even after the OS is wiped!*



SMM vulnerabilities in UEFI

- Privilege escalation from driver execution environment (DXE) to SMM
 - SMM callouts
 - Executing code in SMM that lives outside of protected memory (SMRAM)
 - SMM CommBuffer vulns
 - CommBuffer
 - Type that handles communication between SMM and DXE
 - Copies variables into SMRAM
 - Should check that all nested pointers in a CommBuffer are pointing into SMRAM



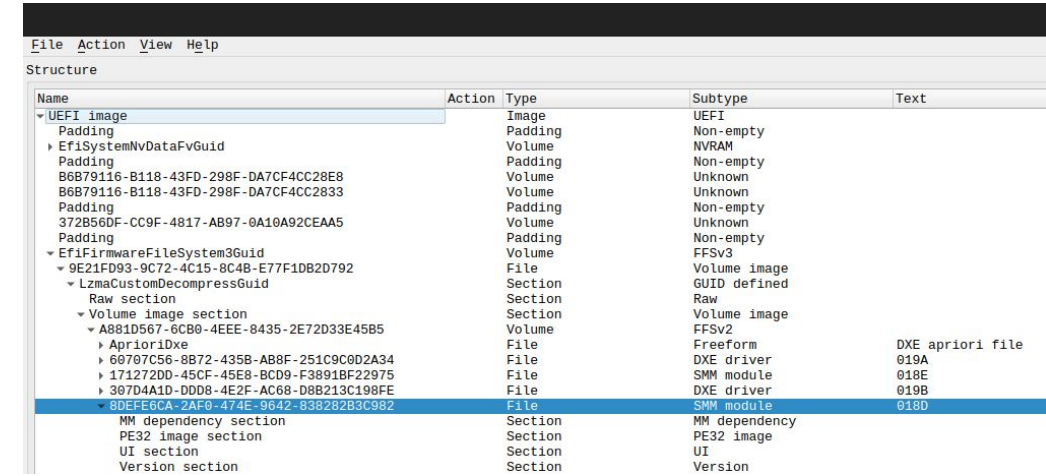
HARDEN

How do we find vulnerabilities in UEFI?

- Manual analysis (slow, requires expert knowledge)
- Fuzz testing (not scalable to an entire UEFI image)
- Alternatively: Use static analysis
 - Trace the flow of data to potentially vulnerable callsites
 - Build dataflow chains
 - Use SMT solving to see if these chains can be exploited
 - (I don't know how to do this part)

Finding SMM vulnerabilities...

- Find vulnerable UEFI drivers
 - Binary writeups
 - Manual, expert analysis of UEFI drivers
 - CVEs
- Download the firmware from the OEM site
- Decompress it with 7-zip (unless you're Dell 😡)
- UEFITool
 - Tool that displays/extracts UEFI drivers in a firmware binary
 - Extract the vulnerable driver, as identified by its GUID



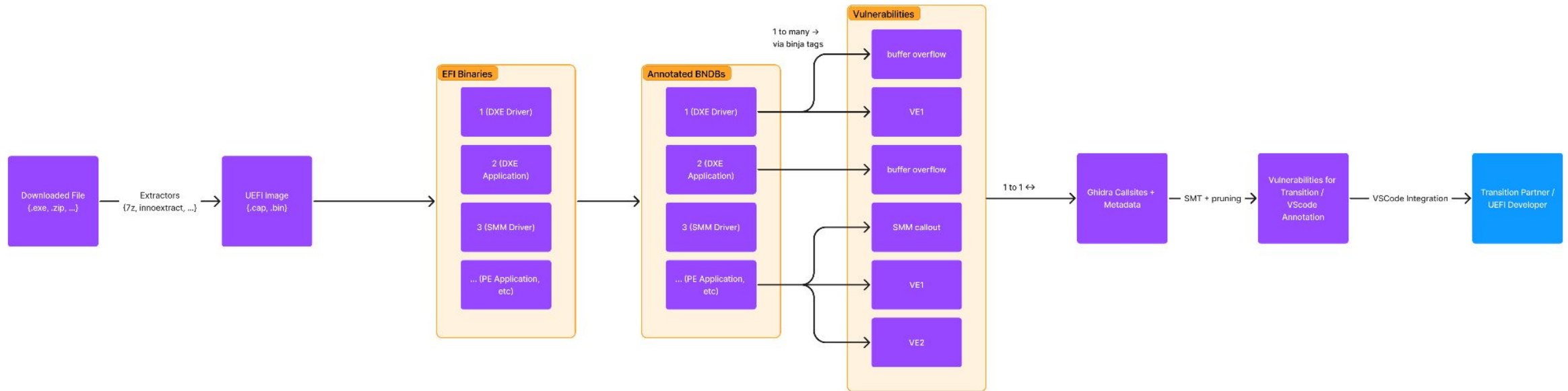
Name	Action	Type	Subtype	Text
UEFI image		Image	UEFI	
Padding		Padding	Non-empty	
EfiSystemNvDataFvGuid		Volume	NVRAM	
Padding		Padding	Non-empty	
B6B79116-B118-43FD-298F-DA7CF4CC28E8		Volume	Unknown	
B6B79116-B118-43FD-298F-DA7CF4CC2833		Volume	Unknown	
Padding		Padding	Non-empty	
372B56DF-CC9F-4817-AB97-0A10A92CEAA5		Volume	Unknown	
Padding		Padding	Non-empty	
EfiFirmwareFileSystem3Guid		Volume	FFSV3	
9E21FD93-9C72-4C15-8C4B-E77F1DB2D792		File	Volume image	
LzmaCustomDecompressGuid		Section	GUID defined	
Raw section		Section	Raw	
Volume image section		Section	Volume image	
A881D567-6CB0-4EEE-8435-2E72D33E45B5		Volume	FFSV2	
AprioriDxe		File	Freeform	DXE apriori file
60707C56-8B72-435B-AB8F-251C9C0D2A34		File	DXE driver	019A
171272DD-45CF-45E8-BCD9-F3891BF22975		File	SMM module	018E
307D4A1D-DDD8-4E2F-AC68-D8B213C198FE		File	DXE driver	019B
80FE6CA-2AF0-474E-9642-838282B3C982		File	SMM module	018D
MM dependency section		Section	MM dependency	
PE32 image section		Section	PE32 image	
UI section		Section	UI	
Version section		Section	Version	

...in Binary Ninja?

- (Arguably) better API with Python support
 - *Jython does not count, Ghidra!*
- More intermediate representations to reason over
 - Single static assignment
- More comprehensive internal tooling
 - PILOT program: already did def/use chaining for vulnerability analysis
- Create a pipeline for automated vulnerability analysis
 - Get binaries, extract them, and run analysis passes over them
 - Use SMT solving to prune the set of possible vulnerabilities



Binary Ninja-based Pipeline



Def/Use Chaining

- Find vulnerable callsites
 - Look for accesses to a particular byte offset from the runtime services table
- From a callsite:
 - Trace the definition of the parameter we want (e.g. *DataSize*)
 - Pointers
 - Trace *down*
 - We don't know how the value the pointer is referencing will change
 - Functions
 - If we know what it does, trace up

```

8 @ 00040cce DataSize @ mem#2 -> mem#3 = 8
9 @ 00040cd6 rax#1, mem#4 = sub_1f6f0(rcx#1, &data_cdffa) @ mem#3
10 @ 00040cde if (rax#1 == 0) then 11 @ 0x40d17 else 17 @ 0x40ce0

11 @ 00040d17 var_38_1#1 = &arg_18
12 @ 00040d1c r9_1#1 = &DataSize
13 @ 00040d20 rax_2#2 = [&RuntimeServices].q @ mem#4
14 @ 00040d27 rdx#1 = &VendorGuid
15 @ 00040d32 rax_3#3, r8_1#1, r9_2#2, mem#5 = rax_2#2->GetVariable @ mem#4(VariableName: u"PlatformLang", VendorGuid: rdx#1, Attributes: 0, DataSize: r9_1#1, Data: var_38_1#1)
16 @ 00040d38 if (rax_3#3 >= EFI_SUCCESS) then 24 @ 0x40d86 else 26 @ 0x40d41

17 @ 00040ce0 rax_1#8 = [&BootServices].q @ mem#4
18 @ 00040ce7 rcx_1#5 = &arg_18
19 @ 00040cf0 mem#11 = rax_1#8->SetMem @ mem#4(Buffer: rcx_1#5, Size: 8, Value: 0) @ mem#4
20 @ 00040cf6 r9#6 = rbx#1
21 @ 00040d02 rcx_2#6 = &arg_18
22 @ 00040d06 mem#12 = sub_211f4(rcx_2#6, &__dos_header_e_cparhdr, &data_122464, r9#6) @ mem#11
23 @ 00040d0b goto 32 @ 0x40d99

24 @ 00040d86 rcx_3#2 = u"PlatformLang reported as %a.\r\n"
25 @ 00040d86 goto 46 @ 0x40d8d

26 @ 00040d41 var_38_2#2 = &arg_18
27 @ 00040d46 r9_3#3 = &DataSize
28 @ 00040d4a rax_4#4 = [&RuntimeServices].q @ mem#5
29 @ 00040d51 rdx_1#2 = &VendorGuid
30 @ 00040d5c rax_5#5, rdx_2#3, r8_1#2, r9_2#4, mem#6 = rax_4#4->GetVariable @ mem#5(VariableName: u"Lang", VendorGuid: rdx_1#2, Attributes: 0, DataSize: r9_3#3, Data: var_38_2#2) @ mem#5 (TraceDirection.UP)

```

```

> 0: rax_5#5, rdx_2#3, r8_1#2, r9_2#4, mem#6 = rax_4#4->GetVariable @ mem#5(u"Lang", rdx_1#2, 0, r9_3#3, var_38_2#2) @ mem#5 (TraceDirection.UP)
> 1: r9_3#3 = &DataSize (TraceDirection.UP)
> 2: r9_1#1 = &DataSize (TraceDirection.DOWN)
> 3: rax_3#3, r8_1#1, r9_2#2, mem#5 = rax_2#2->GetVariable @ mem#4(u"PlatformLang", rdx#1, 0, r9_1#1, var_38_1#1) @ mem#4 (TraceDirection.DOWN)

```

Future Plans

- Static analysis is somewhat limited when it comes to cross-driver interaction
 - Want to evaluate the composability of different vulnerabilities
 - More formal modeling?
 - Emulation?
 - QEMU (OSS)
 - Simics (Intel)
 - More closely replicates the underlying hardware
 - Interrupts between different instructions (which QEMU can't do)
- Leveraging the new tracing engine for SMM vulnerabilities
 - Sort of implemented with Ghidra, but not well
 - Generalize to other CommBuffer vulnerabilities

References & Other Resources

- Rootkits and Bootkits
 - Alex Matrosov (Binarly)
 - Sergey Bratus (DARPA PM for HARDEN)
- [A Tour Beyond the BIOS](#)
 - Jiewen Yao (Intel)
 - Also wrote Securing Firmware (on my reading list)
- SentinelOne blog
- Rylan's excellent JTB from March on UEFI

Acknowledgements

- Jonathan Prokos (mentor for the summer)
- Jacob Denbeaux
- Michael Krasnitski

